

Using Clos Switches in Area Efficient Asynchronous SDM Routers

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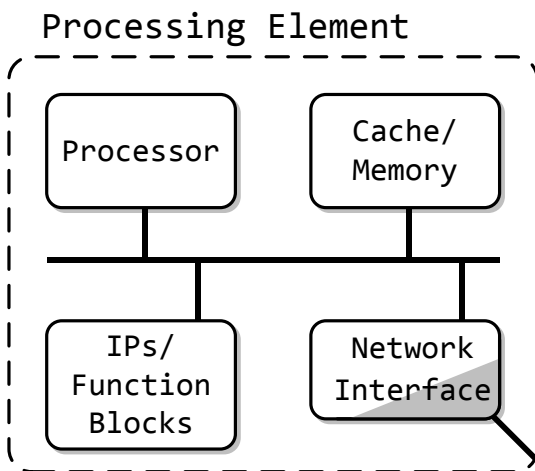
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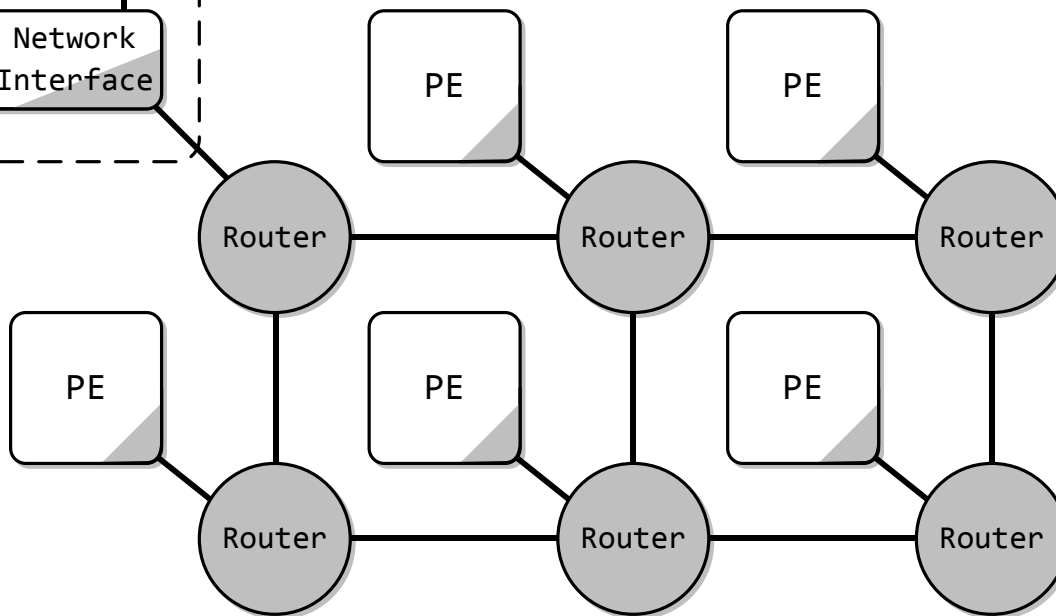
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- Spatial division multiplexing (SDM)
- 2-stage Clos switch
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 - Clos
 - 2-stage Clos switch
- SDM router using Clos switches
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Asynchronous Network-on-Chip

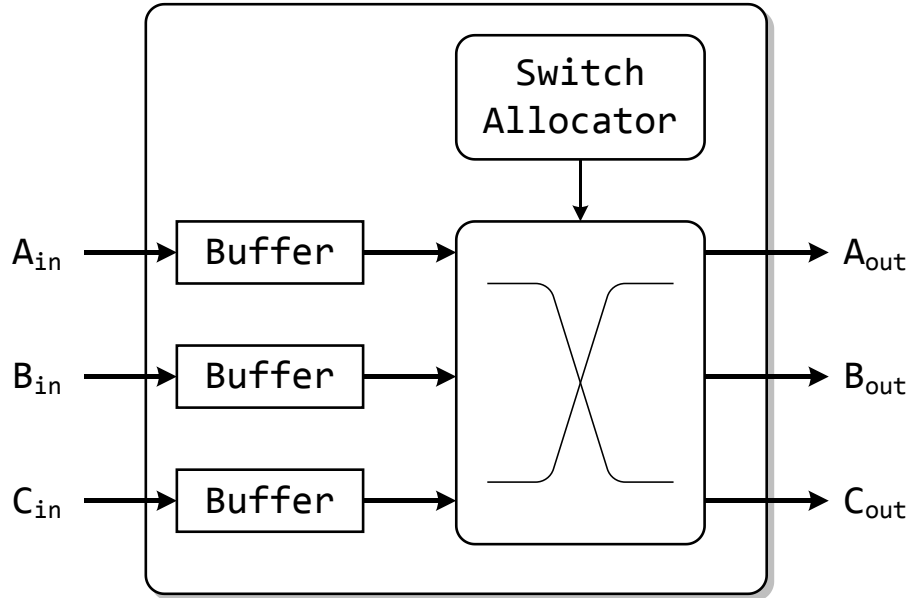


Network-on-Chip (NoC) or **on-chip network** is the state-of-the-art on-chip communication structure for multiprocessor systems.

GALS:
Globally
asynchronous
and locally
synchronous



Basic Router Structure

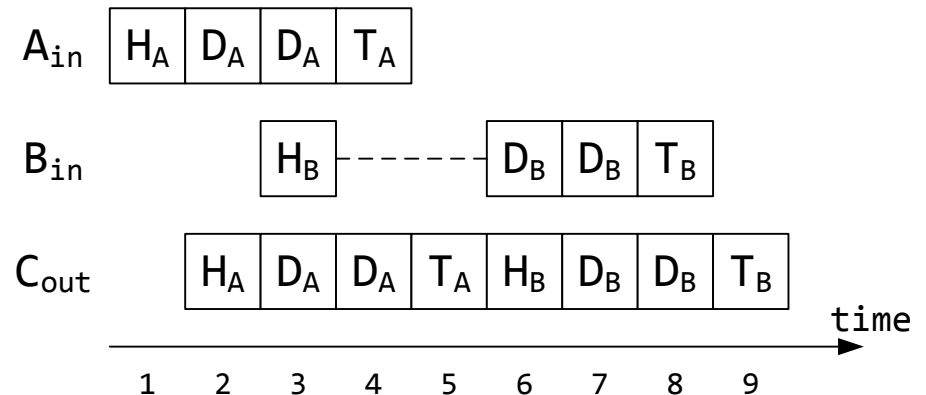


Flow control:

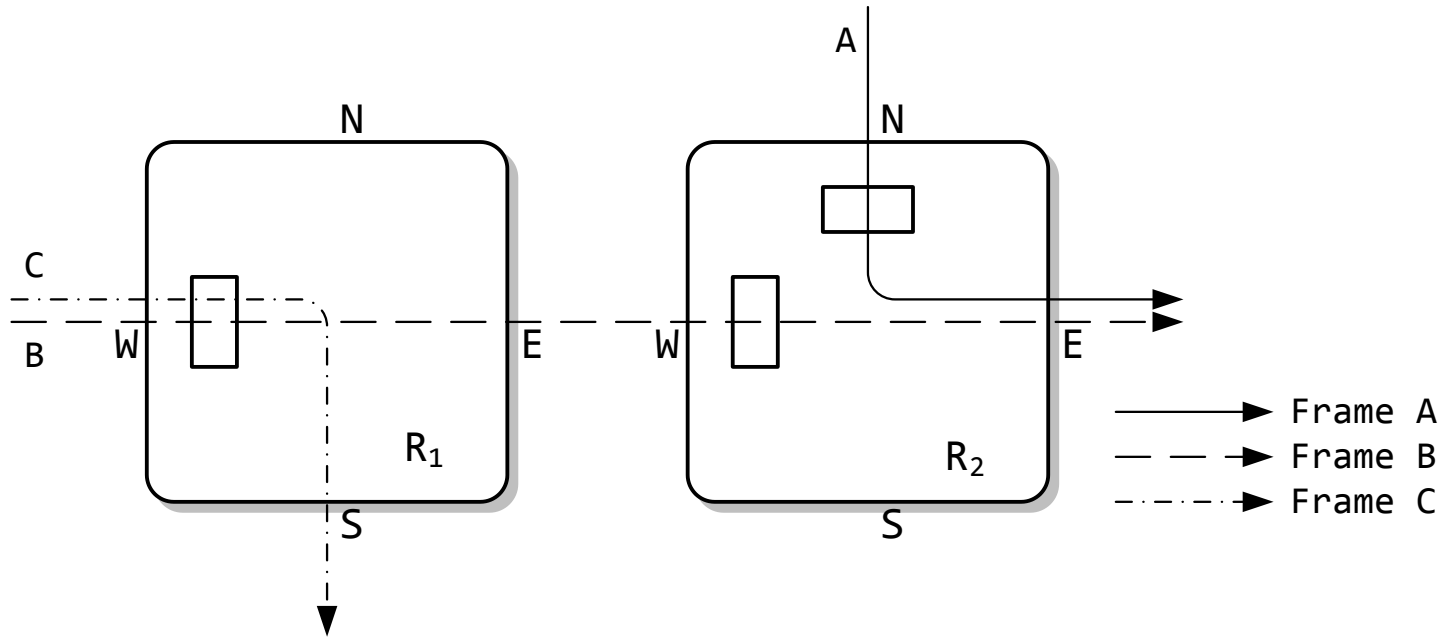
The algorithm used to allocate resources in a router to multiple frames.

Wormhole:

Frames are divided into flits. The header flit contains the target addr. and it is used to reserve a path. Other flits just follow the path.



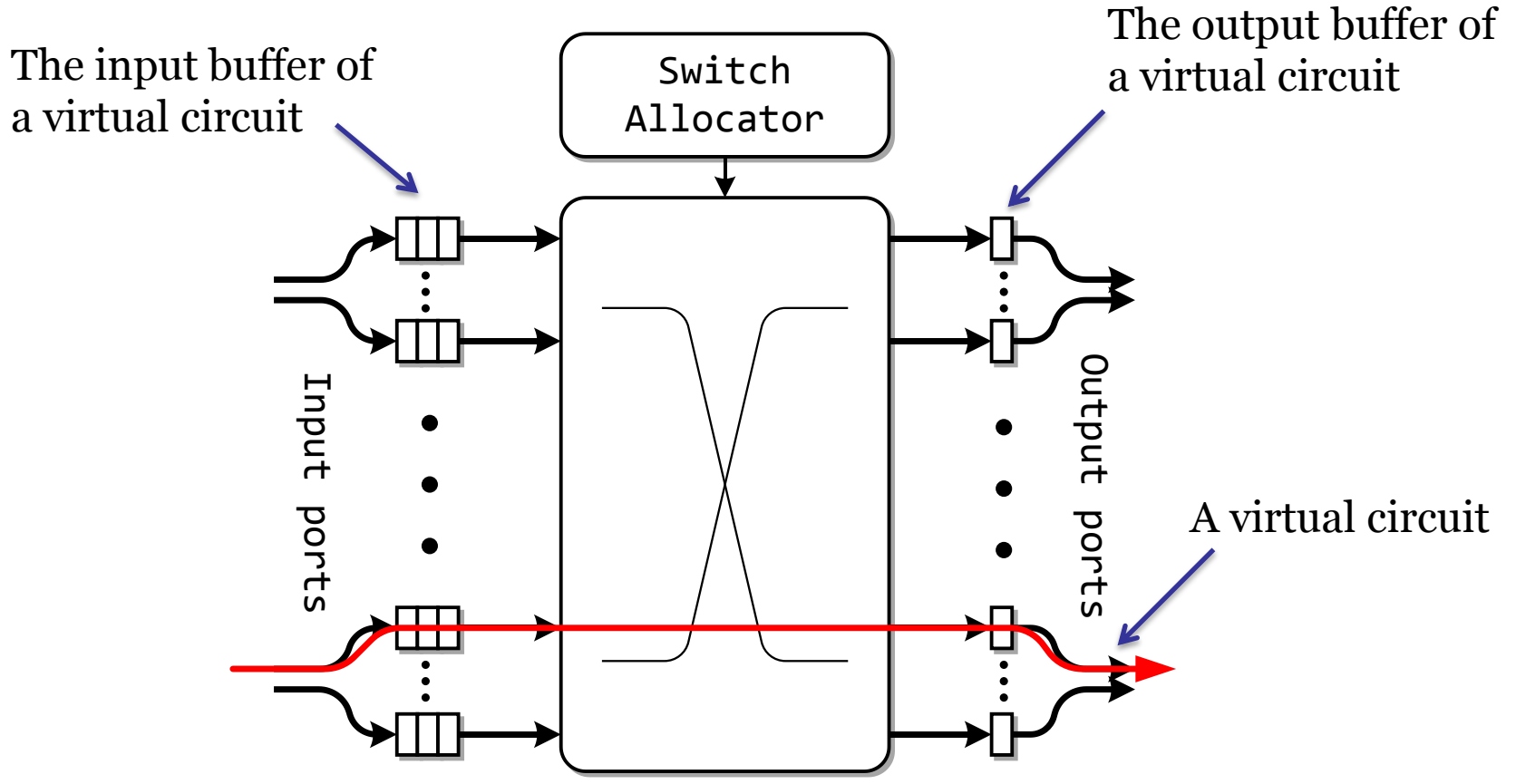
Head-of-Line (HOL)



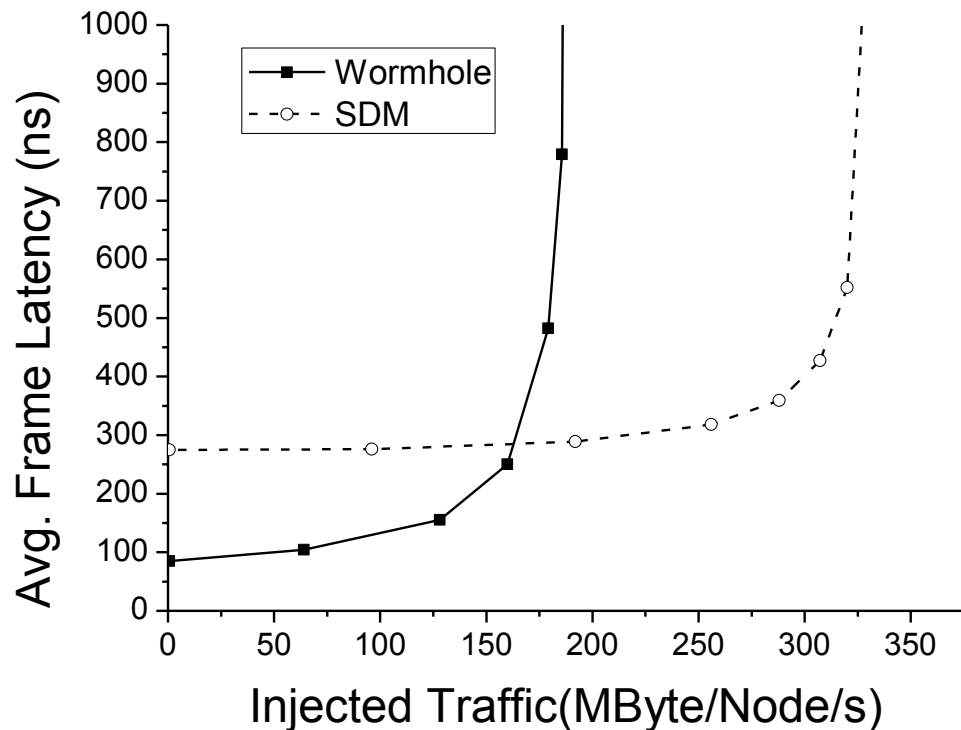
The spread of blocking:

Assuming frame B is blocked by frame A, frame B may also block frame C; however, C is not directly related to A.

Spatial Division Multiplexing



SDM vs. Wormhole



	WH	SDM
Input Buf.	14,303	21,995
Output Buf.	5,935	6,000
Crossbar	4,356	21,744
Arbiters	772	22,208
Overall	25,366	71,956

The area of crossbars

WH:

$$P * P * W$$

SDM:

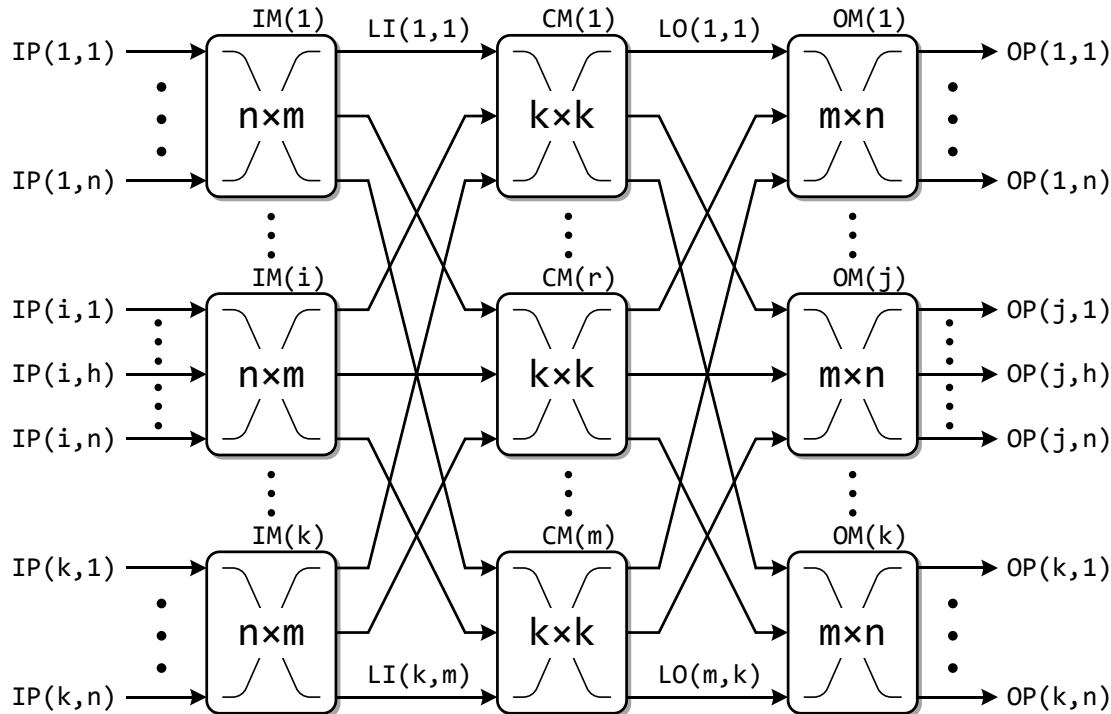
$$MP * MP * W / M = MP * P * W$$

[7] W. Song and D. Edwards. "Asynchronous spatial division multiplexing router," *Microprocessors and Microsystems*, 35(2), 85-97, 2011

Clos Switch - Motivation

- The problems of SDM
 - High-radix crossbars
 - Large crossbar and switch allocator
- Clos networks are the optimal switch structure
- Problems to solve
 - Dynamic configuration [11]
 - Optimal structure for SDM router (this paper)

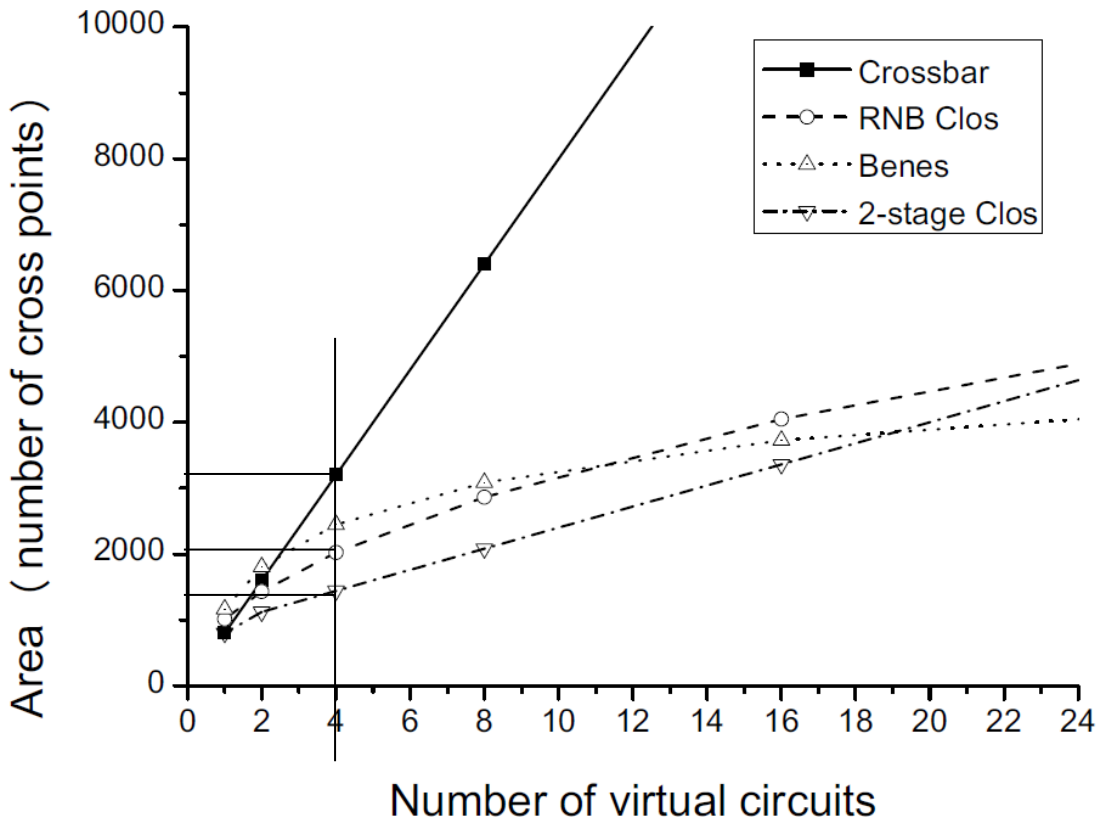
Clos Networks



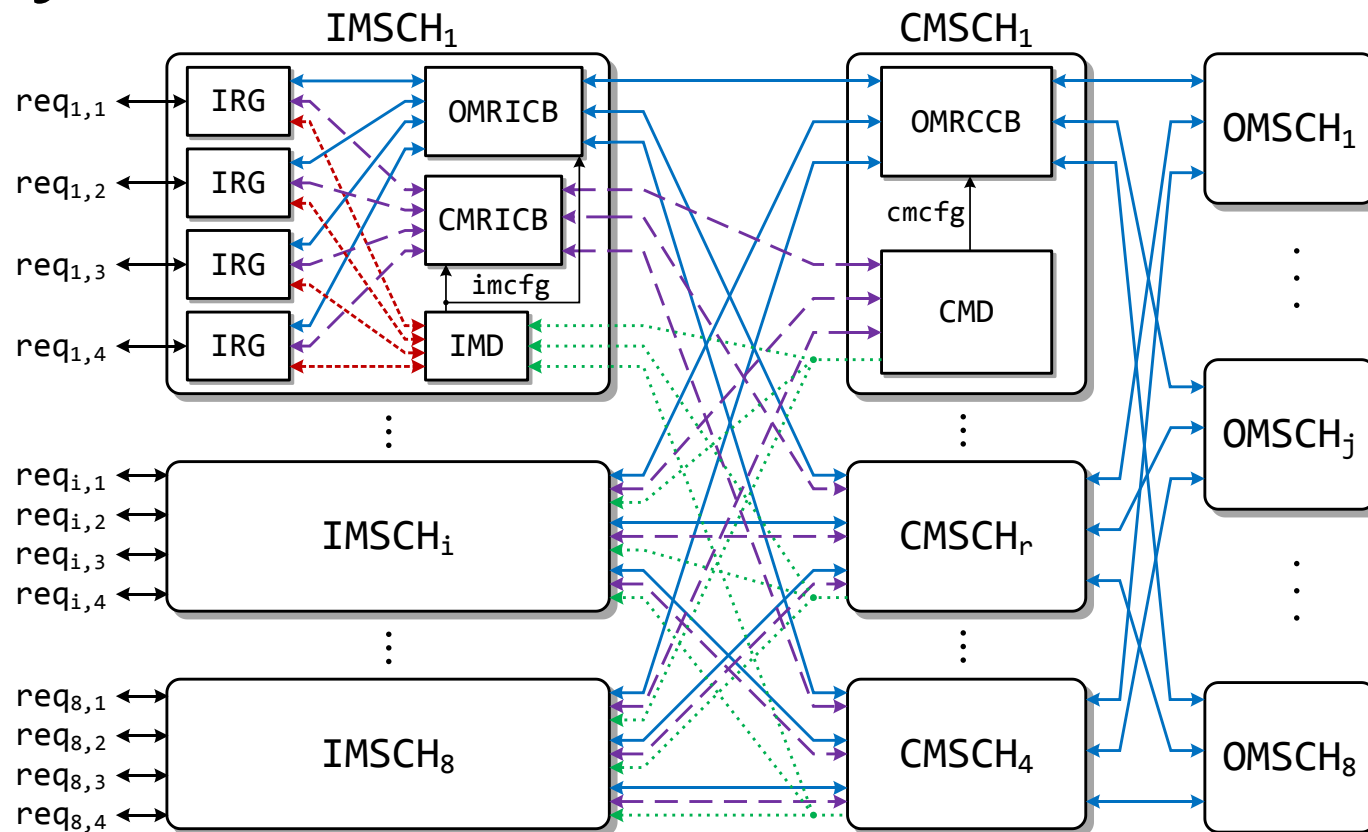
IP/OP: input/output port
IM: input module
CM: central module
OM: output module
 n : number of IPs in IM
 k : number of IMs
 m : number of CMs
 $N = kn$: the total number of IPs

When $m \geq n$, the switch is no-blocking

Clos vs. Crossbar

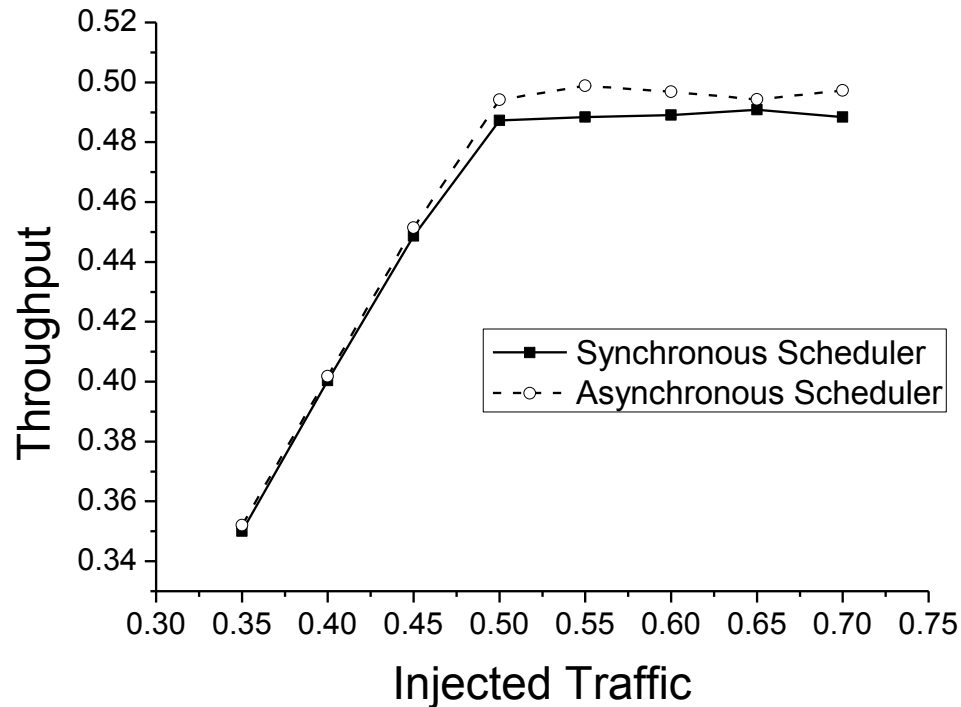


Asynchronous Clos Scheduler



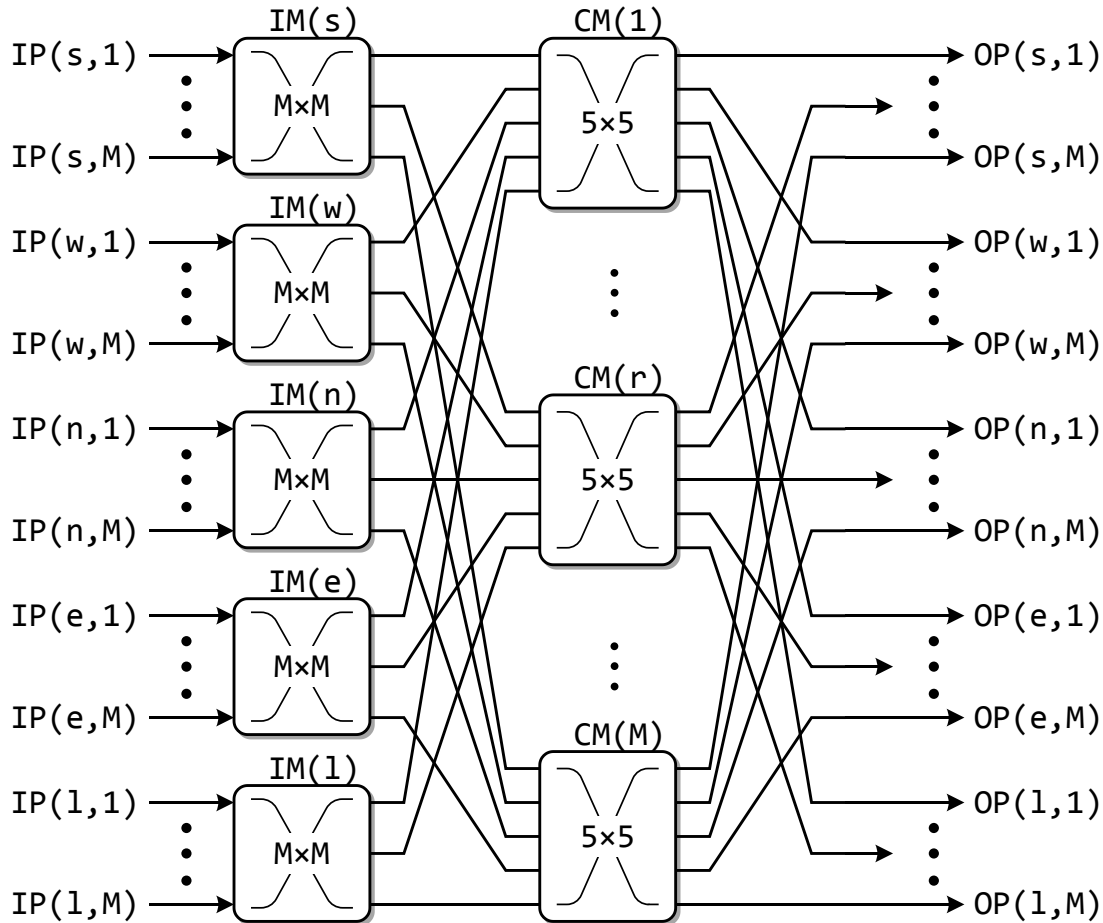
[11] W. Song and D. Edwards. "An asynchronous routing algorithm for Clos networks," In *Proc. of International Conference on Application of Concurrency to System Design*, 2010, 67-76

Async vs. Sync Algorithm



[11] W. Song and D. Edwards. “An asynchronous routing algorithm for Clos networks,” In *Proc. of International Conference on Application of Concurrency to System Design*, 2010, 67-76

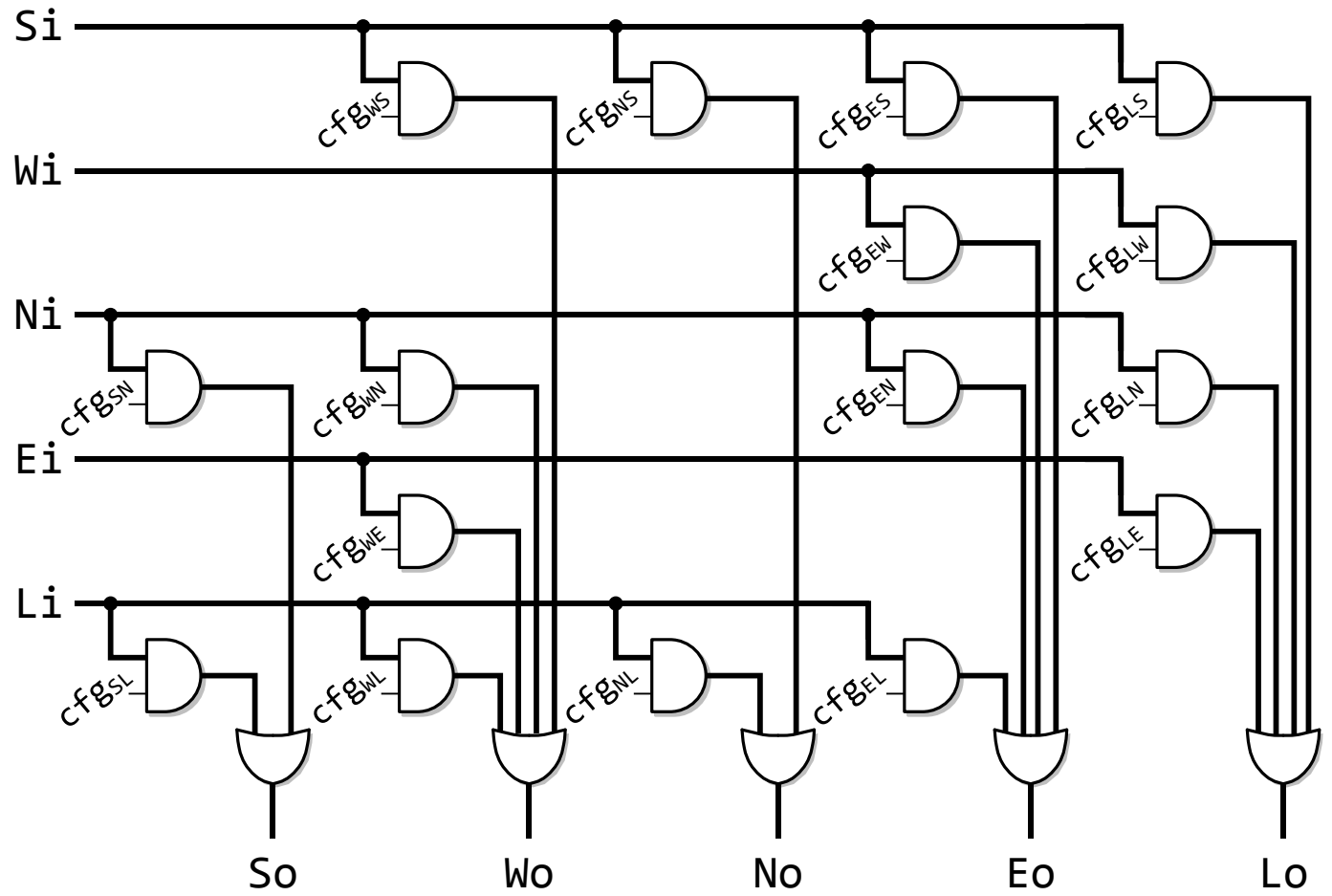
2-stage Clos Switch



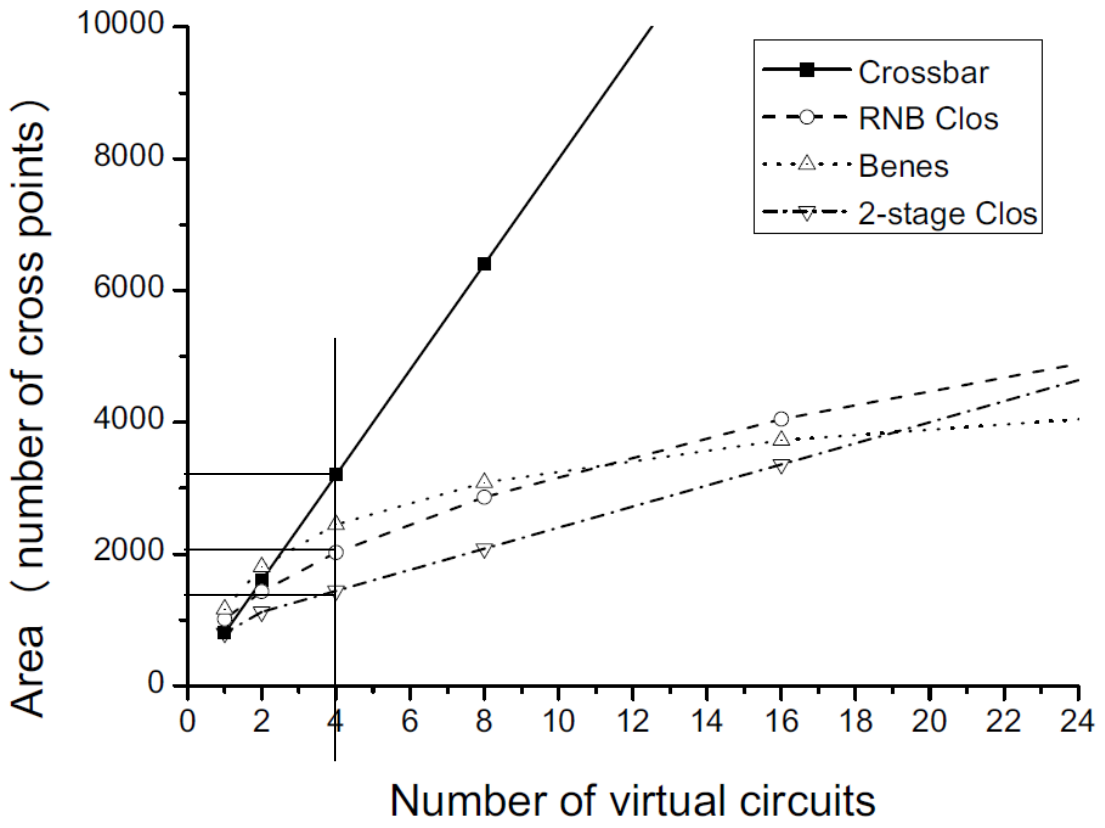
Benefits:

1. 2-stage
 - 1.a less latency
 - 1.b smaller
 - 1.c simpler scheduler
2. CMs can be further simplified

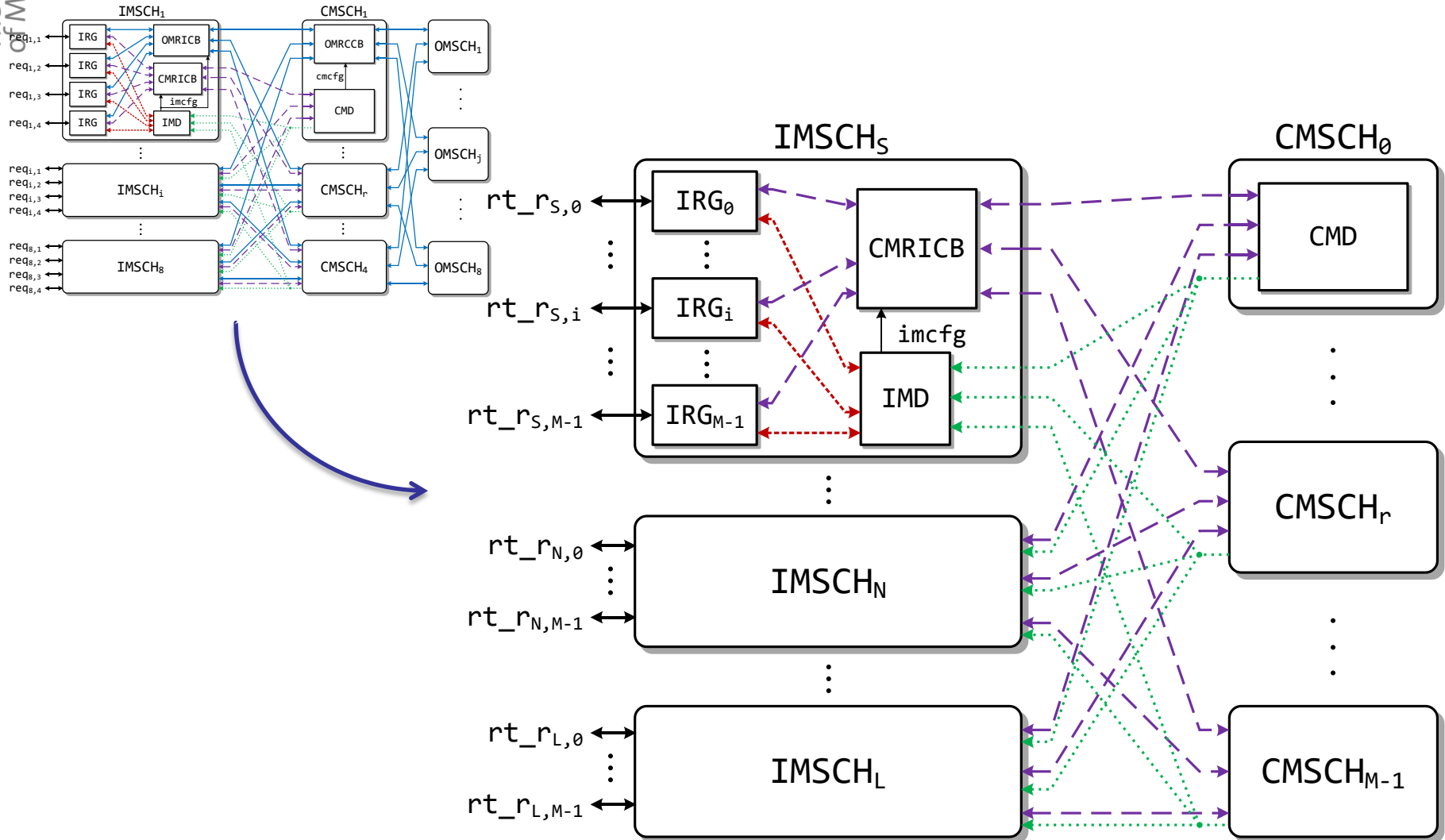
Simplified CM



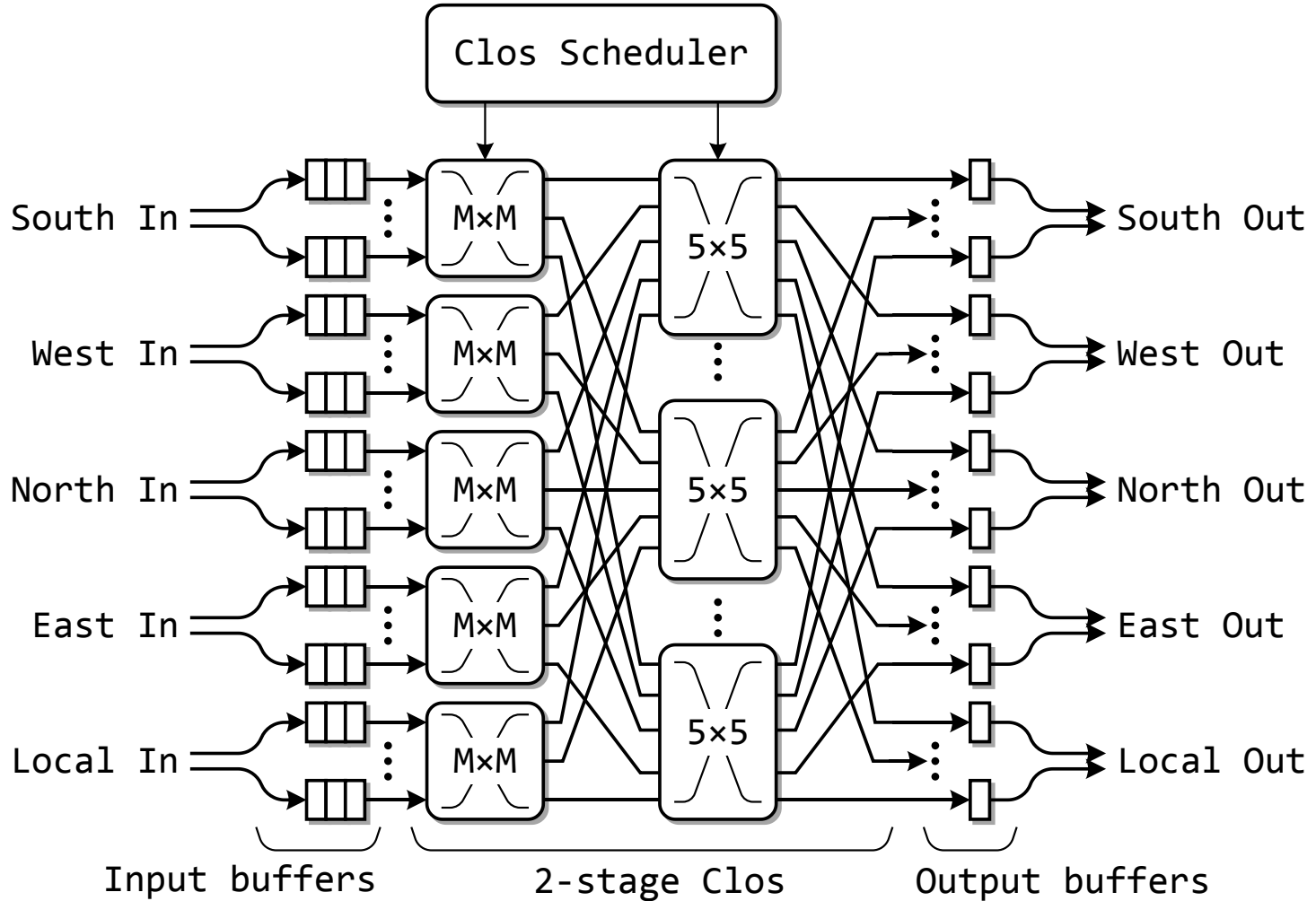
Clos vs. Crossbar



2-stage Clos Scheduler



A New SDM-Clos Router



Area Breakdown

	WH	SDM	SDM-Clos
IP Buf.	17,147	23,356	19,995
OP Buf.	11,034	9,016	8,893
Switch	17,754	69,701	33,417
Alloc.	909	83,485	19,884
Overall	48,170	187,872	84,054

(unit: μm^2)

WH: 1 SDM: 3.9 SDM-Clos: 1.7

Speed Performance

	WH	SDM	SDM-Clos
Cycle period	2.24	2.86	2.67
Router latency	0.92	1.18	1.29
XY router	0.64	0.68	0.63
Allocation	1.24	2.01	2.13

(unit: ns)

Network Performance

	WH	SDM	SDM-Clos
Throughput (MByte/s)	3414.6	3458.1	3458.1
Avg. latency (ns)	108.6	239.4	248.5
Power (mW)	15.6	15.2	13.6

Throughput requirement: ~3,400 Mbyte/s

SDM-Clos:

small latency overhead; low energy consumption; half of the area

Conclusion

- SDM improves throughput
- Clos switch can reduce the area overhead
- A new 2-stage Clos switch
 - Half area (4 virtual circuits)
 - Small latency overhead
 - Less energy consumption

Thanks

http://opencores.org/project,async_sdm_noc